

# The Construction of a Network for Indoor Navigation

Eliseo Clementini<sup>a</sup> and Andrea Pagliaro

*Dept. of Industrial and Information Engineering and Economics, University of L'Aquila, L'Aquila, Italy*

**Keywords:** Indoor Navigation, Network Construction, Straight Skeleton, Visibility Graph.

**Abstract:** Navigation systems provide help to moving agents by enabling them to reach a desired destination. The development of the Global Positioning System (GPS) has enabled the development of outdoor navigation systems, which are based on the knowledge of a map and the user's location and provide guidance indications. The reality of indoor navigation systems is much different: difficulties stem from the fact that the movement of users is not constrained to belong to a well-established network such as the road network in the case of cars, but it is generally freer and less codifiable. This paper focuses on the automatic construction of a navigation graph superimposed on the map of a building that describes the possibilities of movement of a user within the building itself. The construction of the graph tries to replicate the spontaneous movement of users.


## 1 INTRODUCTION

In many situations of real life, it is often necessary to know the path to be performed within a large and complex structure. Navigation systems such as Google Maps or TomTom have from long time addressed such a goal, allowing us to obtain the best route to be followed on foot, by bike or car, based on our GPS position. At the moment, however, navigation systems are limited to exterior use. A current research trend is to expand the navigation technology to inside space, where the GPS signal is not available or its resolution for common handheld devices is not enough. Indoor navigation introduces a methodology that, through the availability of the plans of a building, manages to build a navigation path (Zlatanova et al., 2020; Yang and Worboys, 2015). Navigation implies guidance along the best path to a destination avoiding that the user gets lost (Fallah et al., 2013). Indoor navigation systems have been studied also in reference to building models such as CityGML or IFC (Kolbe et al., 2005).

The focus of this paper is the construction of a network to be overlaid to the building plan for helping navigation. The rationale behind this network is that it should simulate the most natural path of a pedestrian walking inside the building, therefore, avoiding narrow angles, walking not too close to walls and obstacles, taking advantage of visibility of landmarks,

and so on. Much work has already been done for network construction, but the resulting networks always have some critical parts. Taking the medial axis of indoor space, such as in (Yao and Rokne, 1991) and (Lee, 2004), is generally acceptable for long corridors, but it is less natural for large spaces such as halls, where the medial axis loses its significance. In the literature of navigation meshes in simulation and gaming, e.g., (van Toll et al., 2018), medial axis are modified with various techniques to improve path planning. Further, the medial axis of a non-convex shape contains parabolic arcs: those arcs increase the difficulty of computing the medial axis and representations are more complex; further, parabolic arcs do not work well in map-matching algorithms (Taneja et al., 2016). Algorithms based on straight skeletons have been suggested to overcome this limitation (Fu et al., 2020).

Taking the visibility graph between vertices of the shape (Taneja et al., 2011) allows us to consider shorter paths to cross a large area such as a hall, but it leads to unnatural path segments too much close to walls and obstacles. An interesting data structure that modifies a visibility graph with Voronoi diagrams to obtain more natural-looking paths was proposed in (Wein et al., 2007). Hybrid approaches to the determination of navigation networks have been devised to mediate between the above approaches. For example, in (Mortari et al., 2019), a constrained Delaunay triangulation combined with inward offset polygons is able to produce a navigation network that solves the

<sup>a</sup>  <https://orcid.org/0000-0002-3057-7105>

disadvantage of medial axis, but it is critical for doors near vertices of the room, where serpentine routes with too many intermediate points are generated.

In this paper, we conceive a new technique for the network construction, which starts from the straight skeleton of an indoor space. The straight skeleton has no curved segments contrary to the medial axis but it may lead to unnatural narrow turns. Therefore, we smooth the narrow angles of the straight skeleton by avoiding turns bigger than 90°. Then, other points of interest, such as doors, are added to the network. Hence, nodes of the obtained network are shortcut by computing the visibility between nodes, thus suggesting shorter paths. Further simplifications regard the deletion of paths that are too close to walls by intersecting the network with an inner buffer of the indoor space.

In the remainder of this paper, we introduce basic data structures and the straight skeleton algorithm (Section 2). In Section 3, we modify the network to include semantic information such as the doors. In Section 4, we make the network smoother to avoid sharp turns. In Section 5, we add shortcuts between visible points. In Section 6, we remove paths too much closer to the walls. Short conclusions end the paper (Section 7).

## 2 DATA STRUCTURES AND ALGORITHMS

From a cognitive perspective (Montello, 1993), indoor navigation is a problem that can be placed at the environment space, that is, a space that can be perceived by moving and sight, contrary to outdoor navigation, which is placed at geographic scale. Independently of scale, indoor navigation involves the use of methods and algorithms of geographic information science.

An indoor map can be schematized as an irregular tiling, where each polygon models a room or another identifiable portion of space. Even if the rooms are not actually closed, we will consider them as polygons, like the indoor map of Figure 1.

Let us consider the difference between the medial axis and the straight skeleton. Both are algorithms that allow us to extract a representative linear object from a planar shape. The medial axis of a polygon is defined as the place of the points within the polygon which have equal distance from more than one point of the boundary. If a polygon is non-convex, the medial axis has curved parts in correspondence of non-convex vertices of the polygon. To avoid this difficulty, it is possible to consider the straight skele-

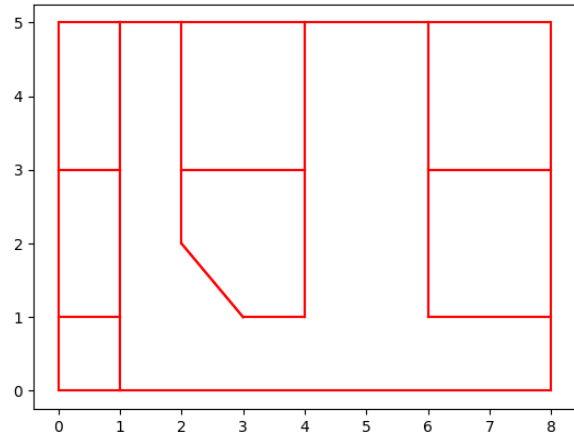


Figure 1: Case study.

ton. The latter is defined as the trace of vertices when a progressive shrinking of the polygon is performed. The two definitions coincide for convex polygons and mark a difference for non-convex polygons, as shown in Figure 2.

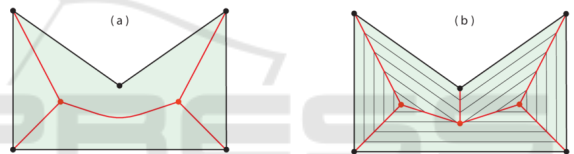


Figure 2: Differences between medial axis (a) and straight skeleton (b).

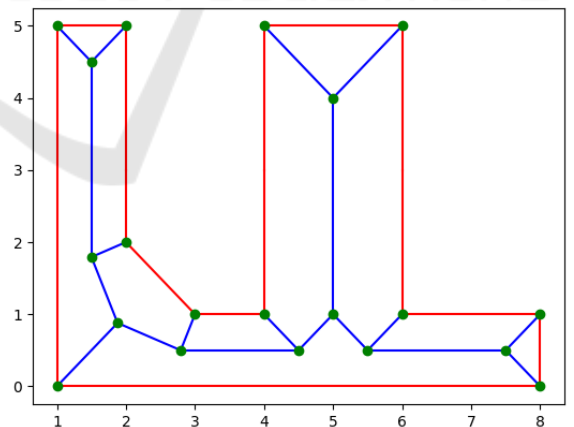


Figure 3: Skeletonized non-convex polygon.

Straight skeletons can be more easily implemented both from the algorithmic point of view and from the necessary data structures. Figure 3 clearly shows the effect after applying the algorithm on a non-convex polygon of our case study map. Storing the results derived from the application of the straight skeleton algorithm requires a suitable data structure, which we

identify as an undirected graph weighed on the Euclidean distance between the nodes, for which we used a Python3 library called “networkx”.

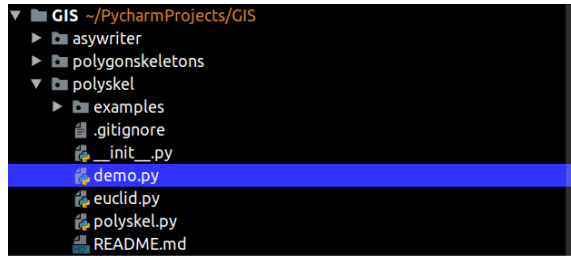


Figure 4: Straight skeleton package organization.

We used an open-source solution of the straight skeleton (github project: <https://github.com/Botffy/polyskel.git>), which represents a Python3 implementation of the algorithm by (Felkel and Obdržálek, 1998). Such a package (Figure 4) is organized as follows:

- demo.py is a main program that imports the points to describe a polygon and its internal holes. This information is enclosed in an external file and imported when it is run.
- skeleton.py is the class that contains all the functions that describe the straight skeleton algorithm.
- euclid.py is a basic Euclidean geometry library that is used to instantiate objects such as points and segments and perform intersection and connection operations between them.

We applied the algorithm to the scenario of Figure 1, after having inserted some holes to simulate the presence of objects inside the rooms and understand how the algorithm behaved in this situation. The algorithm was applied separately to each polygon, obtaining a set of graphs that represent the skeletonized rooms (see Figure 5).

### 3 INSERTING POINTS OF INTEREST

In this section, we describe two aspects: the elimination from the graphs of uninteresting points for navigation, such as the room vertices, and the insertion of semantically rich points, such as the doors.

#### 3.1 Graph Simplification

For navigability, it is not necessary to reach the corners of rooms. So we can eliminate the unnecessary vertices from the graphs by eliminating from each

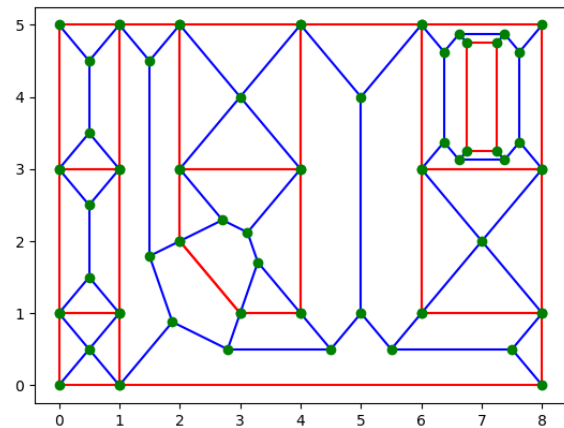


Figure 5: Straight skeletons of all rooms.

graph the nodes with only one incoming arc, that is, the leaves of the graph (Algorithm 1). The result of the simplification is shown in Figure 6, where each room has its simplified graph.

Algorithm 1: Reduce graph of straight skeleton.

```

1: for node in graph.nodes do
2:   if node.degree is < 2 then
3:     delete node from the graph
4:   end if
5: end for
    
```

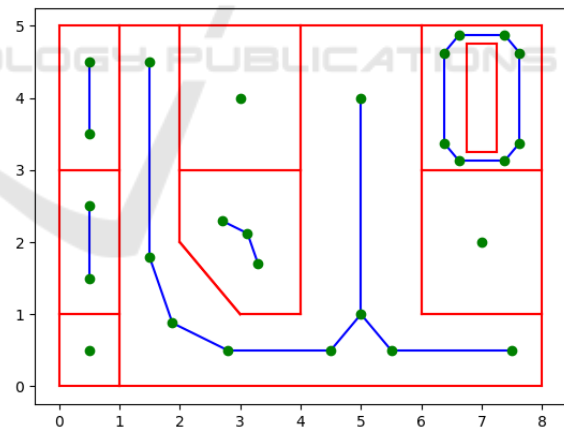


Figure 6: Simplified graphs.

#### 3.2 Inserting Access Points

To this point, we have a set of separate graphs for each room deriving from the straight skeleton. To allow the transition from one room to another, we need to add a connection between the two graphs of adjacent rooms. We build each connection with a central node (defined from now on as a *door node*) connected to the graphs by means of two arcs. Connecting the rooms

is important to apply a search algorithm for the shortest paths between a source and a destination node belonging to different rooms. To facilitate and improve the management of information that allows us to understand which rooms are connected via that door, we define a matrix called the *room-door connection matrix*. This matrix is composed by a door node for each row by and a room for each column.

Algorithm 2: Construction of room-door connection matrix.

```

1: matrix=[]
2: for door_node in door_nodes do
3:   row=[]
4:   for polygon_room in polygon_rooms do
5:     if room intersects door then
6:       add true in row
7:     else
8:       add False in row
9:     end if
10:  end for
11:  add row to matrix
12: end for

```

Starting from a list of door nodes, we calculate if a door intersects a given room and populate the matrix row by row (Algorithm 2). Therefore, in each row of the matrix, there will no more than two matches, since the only cases that can occur are those in which a door belongs to a single room (external door) or two rooms (internal door). On the other hand, in each column, we can observe the number of doors connecting a room.

Algorithm 3: Construction of the door paths.

```

1: door_paths=[ ]
2: for graph room that matches adjacency matrix do
3:   conn distance edges vector
4:   for edge in graph do
5:     find the minimum distance from the door and
     the edge and add to conn
6:   end for
7:   take the edge with minimum distance from conn
     and build a new edge that connects this one with the
     door
8:   if vertex u or v of edge are not in graph then split
     edge
9:   end if
10:  the result is added to door_paths
11: end for

```

Algorithm 3 allows us to obtain the door paths, that is, the pairs of arcs that have doors as the central node and that connect graphs of two different rooms. This algorithm creates the arcs that connect a door node to graph rooms. For each door, the algorithm looks for the closest point between all the arcs of the matching graph. If the closest point is a node of the graph, the

connection is simpler. If the closest point falls inside an arc, then it is necessary to split the arc into two shorter arcs. The results of the application of the operations just mentioned are visible in Figure 7.

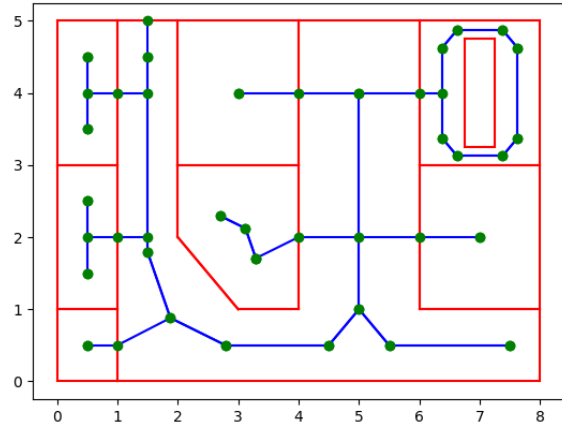


Figure 7: Graph with door paths.

### 3.3 Association of Nodes to Rooms

When joining all the single graphs with the paths on the newly calculated door nodes, it is necessary to keep trace of the room every node belongs to. The door nodes are already represented in the previously mentioned matrix, where they belong to two rooms. For the nodes inside the rooms, Algorithm 4 iterates over each graph and check the nodes by attributing to each of them the identification number of its room.

Algorithm 4: Function to add information about node membership.

```

1: for graph in graphs do
2:   for node in graph.nodes do
3:     add in node info about the membership room
4:   end for
5: end for

```

## 4 SMOOTHING SHARP ANGLES

The straight skeleton algorithm for polygons with internal angles greater than 270° leads to linear shapes with sharp turning angles, like in the case of Figure 8. It is for this kind of situations that the straight skeleton algorithm deviates from the medial axis. In this section, we adopt a strategy of “cutting” the vertices in correspondence of internal angles greater than 270°.

Algorithm 5 identifies the vertices to be cut. Each identified vertex is substituted in the room boundary by two vertices at a small distance. In Figure 8 and

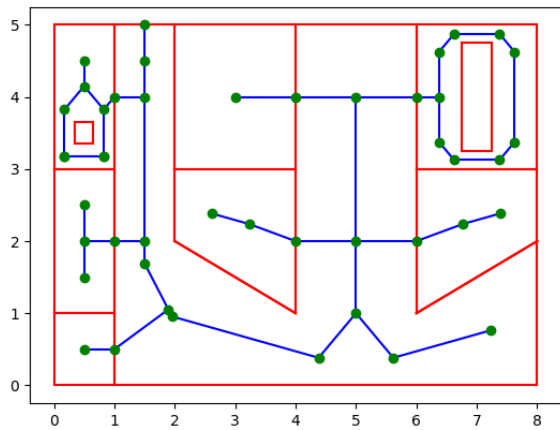


Figure 8: The network before the application of function to smooth sharp vertices.

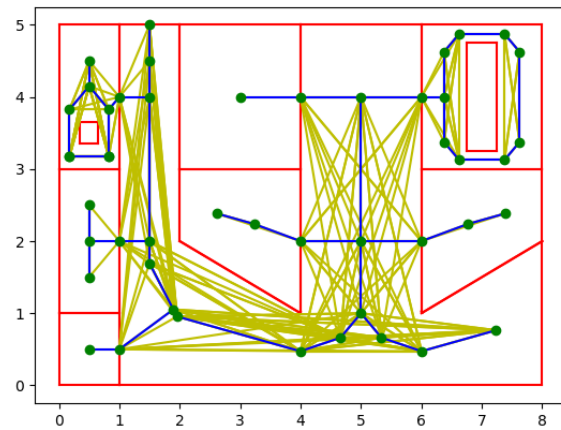


Figure 10: Overlay with visibility graph.

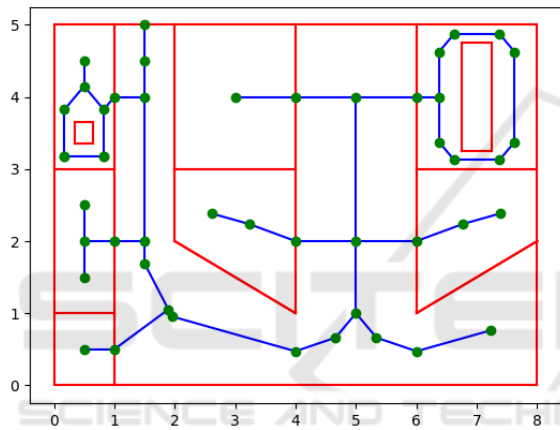


Figure 9: The network after the application of function to smooth sharp vertices.

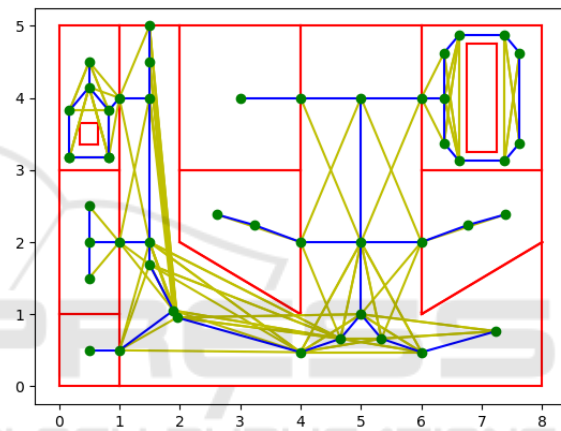


Figure 11: Visibility graph after removal of intersections.

Algorithm 5: Function to smooth angle.

```

1: procedure AnglePolygonSmooth(rooms,c)
2:   for room in rooms do
3:     for poly in room do
4:       for incremental triple of vertex do
5:         compute the angle between them
6:         if internal angle is > 270 degrees then
7:           needs to be cut
8:         end if
9:       end for
10:    end for
11: end procedure
    
```

Figure 9, it is possible to compare the map before and after the application of this algorithm, with the choice of a very small factor for cutting the sharp vertex. In such a figure, it is difficult to notice the removal of the vertex and its substitution with two extra vertices, but the change in the straight skeleton is evident and qualitatively more appealing for navigation.

## 5 ADDING VISIBILITY GRAPH

As explained in the introduction, in our strategy we combine straight skeletons with visibility graphs. By connecting nodes that are in clear visibility, we obtain shorter paths that mimic the behaviour of other categories of navigation algorithms. Therefore, we use the network nodes already found in previous section and we compute the visibility graph based on these nodes. Algorithm 6 allows the creation of the visibility graph (see Figure 10).

After the visibility graph is obtained, we reduce the number of arcs by eliminating the arcs that intersect with existing arcs of the initial network (see Figure 11). The intersection between a point and a line necessary in Algorithm 6 was obtained with Algorithm 7. The latter calculates the collinearity between a point and a segment via a cross product.

Algorithm 6: Function that checks the visibility between two nodes.

```

1: procedure visibility_check(first_node,second_node,
   room)
2:   intersection = False
3:   for wall in room do
4:     if wall intersects segment composed by first
       node and second node then
5:       intersection = True
6:     end if
7:   end for
8:   if intersection is False then append edge to the vis-
       ibility graph
9:   end if
10: end procedure

```

Algorithm 7: Function to compute intersection between Point and Segment.

```

1: function IntersectionBetweenPointandSegment(P,S)
2:    $M = \begin{bmatrix} S.p1.x & S.p1.y & 1 \\ S.p2.x & S.p2.y & 1 \\ P.x & P.y & 1 \end{bmatrix}$ 
3:   if  $\det(M) = 0$  then ▷ collinearity test
4:      $D(S.p1, S.p2)$  ▷ distance between S.p1 and
       S.p2
5:     if  $D(S.p1, P) < D(S.p1, S.p2)$  and  $D(S.p2, P) <$ 
        $D(S.p1, S.p2)$  then
6:       return True
7:     else
8:       return False
9:     end if
10:  else
11:    return False
12:  end if
13: end function

```

## 6 NETWORK BUFFERIZATION

A problem of a navigation network based on visibility is that produces links that are too much close to walls. As a final phase of the network construction, we avoid such a problem by applying an inward buffer to the polygons that make up the building plan and then removing the links that intersect the buffer.

The width of the buffer can be chosen depending on the extension of the room: therefore, we make it dependent from the apothem of each polygon. The implementation of the buffer for our study case has been made by translating each segment representing a wall by the chosen width and connecting them by circular arcs centered on segment endpoints. Then, a method has been added to verify the intersection with navigation network. The doors between rooms have not been bufferized. The application of the method just discussed on the visibility graph is visualized in Figure 12.

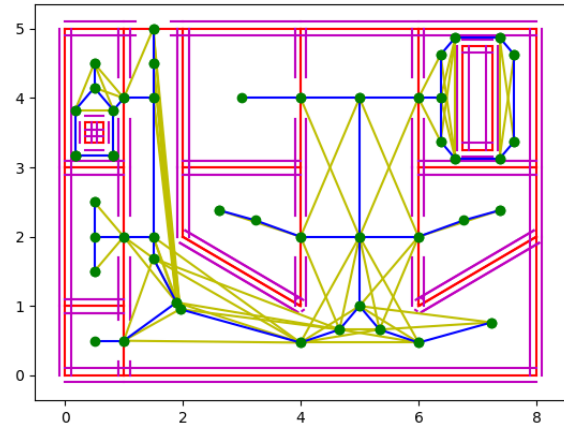


Figure 12: Bufferization of the rooms.

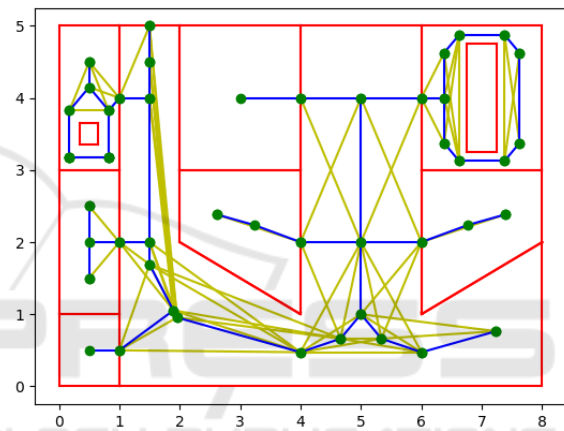


Figure 13: Final graph.

Algorithm 8: Function to filter not connected elements.

```

1: procedure FilterNotConnected(principal_node)
2:   for node in graph.nodes do
3:     if Dijkstra from principal_node to node outputs
       not achievable then
4:       delete node
5:     end if
6:   end for
7: end procedure

```

It can be seen that in this graph there are several eliminations of previous arcs. During this process, it can happen that the elimination of some arcs caused the disconnection of graph components making the graph no longer connected. To verify the reachability of all the nodes of the graph, we choose a main node (which could be for example the entrance of the building) and apply Dijkstra from the main node to other nodes. If reachability is not verified, those nodes will be eliminated making the graph connected again. This function is represented in Algorithm 8. Normally, choosing a thin buffer does not provoke the interrup-

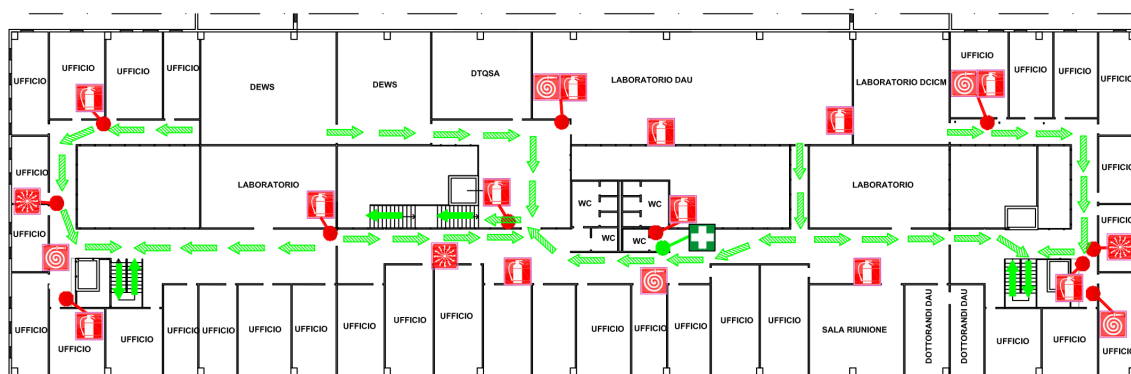


Figure 14: Emergency plan of university floor.

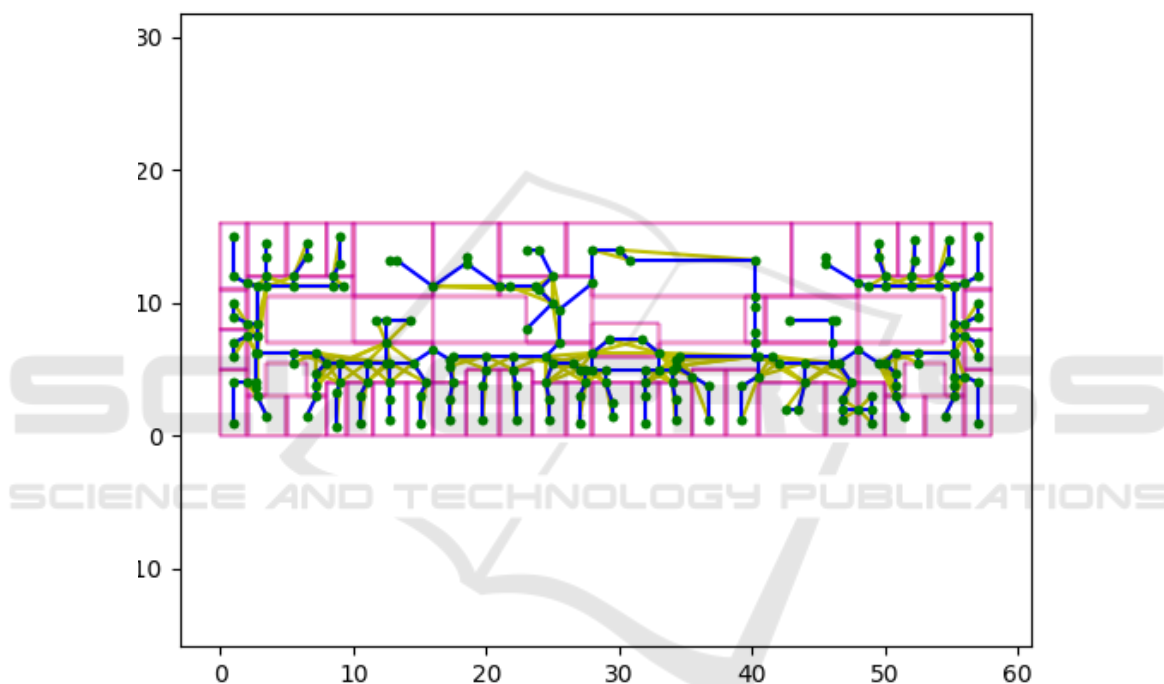


Figure 15: Navigation network for university floor.

tion of the navigation graph. Therefore, this action can be useful if from the main entrance the navigating agent has a considerable size and we want to discard the parts of the map where the agent does not fit. In Figure 13, there is the final navigation network for the case study.

## 7 CONCLUSIONS

In this paper, we proposed the automatic construction of a navigation network for the map of a typical building floor. While there are several proposals in the literature (e.g., (Lee, 2004; Taneja et al., 2011; Mortari

et al., 2019)), there are recognized drawbacks in each proposal. In our contribution, we try to overcome all the drawbacks. We start from computing the straight skeleton of the map and combine it with a visibility graph over the straight skeleton nodes. We believe that this approach interprets a good qualitative modeling of the movement of people inside an indoor environment, obtaining a smooth path at the center of corridors and around vertices and finding shorter paths on larger halls based on visibility. Semantic aspects have been introduced to keep trace of rooms and doors connecting them. An inward bufferization can also avoid movements too close to walls and solve navigation planning for sizeable agents (e.g., for people with reduced mobility).

At the time of writing, we started an evaluation of the algorithm for some real building floors. In Figure 14, we took into consideration an emergency map of a floor in our university and we applied the algorithm to find the navigation network for this map: the result is shown in Figure 15. Further work will follow in several research directions. To validate the shape of the network, we plan to collect various trajectories of moving people inside a building instructing them to walk from the entrance to a given target and compare their trajectories with the proposed navigation network to assess if it could be considered as the representation of the average trajectory. Another development is to find an automatic way of extracting qualitative directions for moving inside the building floor, similarly to the work of (Russo et al., 2014). Directions should not be expressed in terms of angles and metric distances, but in qualitative terms, making use of various models for qualitative spatial reasoning (e.g., (Clementini, 2013; Fogliaroni and Clementini, 2015; Bartie et al., 2013; Clementini and Cohn, 2014; Tarquini and Clementini, 2008)). An extension of the proposed network is necessary as well to connect building floors among them via stairs or elevators and with outdoor space.

## REFERENCES

- Bartie, P., Clementini, E., and Reitsma, F. (2013). A qualitative model for describing the arrangement of visible cityscape objects from an egocentric viewpoint. *Computers, Environment and Urban Systems*, 38(1):21–34.
- Clementini, E. (2013). Directional relations and frames of reference. *GeoInformatica*, 17(2):235–255.
- Clementini, E. and Cohn, A. (2014). RCC\*-9 and CBM\*. In Duckham, M., Pebesma, E., Stewart, K., and Frank, A., editors, *GIScience 2014: Geographic Information Science*, volume 8728 of *Lecture Notes in Computer Science*, pages 349–365, Cham. Springer.
- Fallah, N., Apostolopoulos, I., Bekris, K., and Folmer, E. (2013). Indoor human navigation systems: A survey. *Interacting with Computers*, 25(1):21–33.
- Felkel, P. and Obdržálek, S. (1998). Straight skeleton implementation. In Szirmay-Kalos, L., editor, *Proceedings of Spring Conference on Computer Graphics*, pages 210–218, Budmerice, Slovakia.
- Fogliaroni, P. and Clementini, E. (2015). Modeling visibility in 3D space: A qualitative frame of reference. In Breunig, M., Al-Doori, M., Butwilowski, E., Kuper, P., Benner, J., and Haeefe, K., editors, *9th International 3DGeoInfo Conference, 2014*, Lecture Notes in Geoinformation and Cartography, pages 243–258, Cham. Springer.
- Fu, M., Liu, R., Qi, B., and Issa, R. R. (2020). Generating straight skeleton-based navigation networks with industry foundation classes for indoor way-finding. *Automation in Construction*, 112:103057.
- Kolbe, T., Gröger, G., and Plümer, L. (2005). CityGML: Interoperable access to 3D city models. In van Oosterom, P., Zlatanova, S., and Fendel, E., editors, *Geoinformation for Disaster Management*, pages 883–898, Berlin, Heidelberg. Springer.
- Lee, J. (2004). A spatial access-oriented implementation of a 3D GIS topological data model for urban entities. *GeoInformatica*, 8(3):237–264.
- Montello, D. R. (1993). Scale and multiple psychologies of space. In Frank, A. U. and Campari, I., editors, *Spatial Information Theory A Theoretical Basis for GIS*, pages 312–321, Berlin, Heidelberg. Springer.
- Mortari, F., Clementini, E., Zlatanova, S., and Liu, L. (2019). An indoor navigation model and its network extraction. *Applied Geomatics*, 11(4):413–427.
- Russo, D., Zlatanova, S., and Clementini, E. (2014). Route directions generation using visible landmarks. In *6th ACM SIGSPATIAL International Workshop on Indoor Spatial Awareness, ISA 2014*, pages 1–8, New York, NY, USA. Association for Computing Machinery.
- Taneja, S., Akinci, B., Garrett, J. H., and Soibelman, L. (2016). Algorithms for automated generation of navigation models from building information models to support indoor map-matching. *Automation in Construction*, 61:24–41.
- Taneja, S., Akinci, B., Garrett, J. H., Soibelman, L., and East, B. (2011). Transforming IFC-based building layout information into a geometric topology network for indoor navigation assistance. In *International Workshop on Computing in Civil Engineering 2011*, pages 315–322. American Society of Civil Engineers (ASCE).
- Tarquini, F. and Clementini, E. (2008). Spatial relations between classes as integrity constraints. *Transactions in GIS*, 12(SUPPL. 1):45–57.
- van Toll, W., Cook Iv, A. F., van Kreveld, M. J., and Geraerts, R. (2018). The medial axis of a multi-layered environment and its application as a navigation mesh. *ACM Transactions on Spatial Algorithms and Systems*, 4(1):1–21.
- Wein, R., van den Berg, J. P., and Halperin, D. (2007). The visibility-Voronoi complex and its applications. *Computational Geometry*, 36(1):66–87.
- Yang, L. and Worboys, M. (2015). Generation of navigation graphs for indoor space. *International Journal of Geographical Information Science*, 29(10):210–218.
- Yao, C. and Rokne, J. (1991). A straightforward algorithm for computing the medial axis of a simple polygon. *International Journal of Computer Mathematics*, 39(1-2):51–60.
- Zlatanova, S., Yan, J., Wang, Y., Diakité, A., Isikdag, U., Sithole, G., and Barton, J. (2020). Spaces in spatial science and urban applications—state of the art review. *ISPRS International Journal of Geo-Information*, 9(1):58.