# A HYBRID CONTROLLER FOR A NONHOLONOMIC CAR-LIKE ROBOT

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Hybrid Systems, Switched Systems, Finite Horizon Control, Supervisory Control, Nonholonomic robot. Keywords:

Abstract:

This paper presents the usage of hybrid systems to develop an adaptive recent horizon controller for a nonholonomic car-like robot. The system is modeled by a non-deterministic hybrid system in which transitions represent the discrete control actions, mode invariants the constraints and the transition relation encodes sequencing requirements. The control algorithm examines at runtime the possible traces into the future by determining at which time point to switch to which mode. Based on these predictions the next control move is performed. We demonstrate our approach by controlling a car-like robot though a maze without any pre-runtime path planing.

#### INTRODUCTION

In many applications a high level digital controller is interacting with a set of low level controllers and actuators. Often these low level controllers and actuators, being for example PID controllers, accept only a finite set of discrete control actions, e.g. start PID controller i, turn left, open valve. While there is a finite number of possible control actions, the control input itself is continuous. For a car the actions are turn steering wheel left, turn it right or hold it at its current position, accelerate, hold speed, de-accelerate. Clearly, there are constraints, the steering wheel can only turned right until it reaches its maximal right position, acceleration and de-acceleration rates are finite.

Such control systems can perfectly be modeled using hybrid systems (Almeida et al., 1997; Abdelwahed et al., 2002; Kowalewski et al., 2000; v. Mohrenschildt, 2001; Abdelwahed et al., 2002; Labinaz et al., 1997; Bemporad et al., 2000; Bemporad and Morari, 1999). A hybrid system has continuous states that evolve according to differential equations, and a finite set of modes that evolve according to transitions. Changing mode causes a discrete change of the differential equations. The transitions in between modes are be guarded by conditions on the state variables. Further, each mode contains a mode invariant, a predicate on the state variables that has to remain true while the system is in that mode. These hybrid systems are non-deterministic in their transitions, a transition can occur if its guard allows it, but some transition has to occur if the invariant in a mode becomes false. The goal of the control algorithm is to trigger the transitions, at runtime, such that the system exhibits the desired behavior; minimizing some cost function J. The controller proposed in this paper is adaptive in the sense that the control frequency, the amount of time in between control moves, is not fixed but changes. If becomes shorter as there are constraints and longer otherwise. The control algorithm bases its decision on predictions of the future behavior of the system. As found in many systems in nature, a human driving a car, the control algorithm sees the constraints, being walls or bends in the road only within a finite prediction horizon.

The approaches presented in (Almeida et al., 1997), the authors develop a hybrid controller to control a non-holonomic vehicle that inspired us to apply our approach to motion control problems, and (Abdelwahed et al., 2002; Kowalewski et al., 2000) developing a controller for a two tank system, relaying on apriory analysis of the control problem to develop there hybrid systems. In contrast, the constructions of our controller is generic, resulting in a non-deterministic hybrid system, only the prediction horizon has to be determined by analyzing the system to be controlled. The transition conditions that define the control strategy are determined at run time by the supervisor algorithm.

Several questions arise: Can the control algorithm

prevent the system from entering a blocking state, does the control algorithm control the system into its global minimum with respect to the cost function or do we get stuck in some local minimum. We show that under some general assumptions the proposed control algorithm can satisfy both of these goals. The approach taken follows the ideas presented in (Mayne and Michalska, 1990; Primbs and Nevistic, 2000; Almir and Bornard, 1995) by choosing a prediction horizon such that the finite horizon controller is stabilizing.

We apply the proposed approach to the problem of controlling a non-holonomic car-like robot. The goal of the algorithm is to control a robot though a maze. The algorithm does not have a map of the maze, neither is there any pre-runtime motion planing. The algorithm can only explore the maze within a limited horizon of its current position. We show simulations using a tracked vehicle, a vehicle that can turn on a dime, and a car-like wheeled vehicle. The simulations where performed using a generic implementation of the developed control algorithm in conjunction with C-code that was automatically generated using the computer algebra system Maple from the symbolic solutions of the differential equations.

First we introduce the hybrid systems used to model the system. Section 4 presents the control algorithm and investigates stability conditions. In section 5 we then apply our control algorithm to steer a carlike robot though a maze.

# 2 HYBRID SYSTEMS, EVENTS, CONTROLLED TRACE

A vast amount of literature contains detailed discussion with different approaches to hybrid systems. (Branicky, 1998; Gollu and Varaiya, 1989; Henzinger, 1996; Labinaz et al., 1997).

To motivate the use a non-deterministic hybrid system as a model for control systems we give a simple example of the movements of the steering wheel in a car. The position of the steering wheel is given by the angle  $\phi$ , where  $\frac{-\pi}{4} \leq \phi \leq \frac{\pi}{4}$ . The steering wheel can either be turned left, turned right or held at its position. The turning rate of the steering wheel is limited:  $\dot{\phi} = -0.01, \dot{\phi} = 0, \dot{\phi} = 0.01$ . The hybrid system hence has three modes: TWL,STOP,TWR, turn wheel left, stop, turn wheel right. The turning left and turning right modes have invariants that ensure that the wheel is not turned to far:  $I(\text{TWL}) = \phi \geq \frac{-\pi}{4}$  and  $I(\text{TWR}) = \phi \leq \frac{\pi}{4}$ . In each mode  $\phi$  is computed by a differential equation:  $F(\text{STOP}) := \dot{\phi} = 0$ ,  $F(\text{TWL}) := \dot{\phi} = -0.01$ . The transition guards are shown in the graphical rep-

resentation Figure 1. Note, that if the system is in the TWL mode and  $\phi$  reaches  $\frac{-\pi}{4}$  then the system has a forced transition to the STOP mode. Any trace of this hybrid system represents a "legal" motion of the steering wheel.

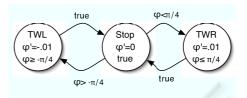


Figure 1: Steering Hybrid System

**Definition 2.1 (Hybrid System)** A hybrid system is a tuple

$$HS = (Q, X, Y, V, F, \delta, I, G, A)$$

where:

- $Q = \{q_0, q_i, \dots, q_e\}$  is the set of discrete states, called modes, of the system,  $X \subseteq \mathbb{R}^{n+m}$  is the set of the continuous states, called states,  $Y \subseteq X, Y \in \mathbb{R}^m$  is the output space.
- V is the set of variables. We distinguish two type of variables, state variables  $x_i, 1 \leq i \leq n$  and output variables  $u_i, 1 \leq i \leq m$ , that are also state variables.
- $F: Q \to (X \to X)$  gives a set of differential and algebraic equations that describe the continuous behavior of the system for each mode.
- $\delta \subseteq (Q \times Q)$  is the transition relation that describes the discrete behavior of the system.
- $I:Q \to (X \to \mathsf{bool})$  gives an invariant for each mode.
- $G:(Q\times Q)\to (X\to \mathsf{bool})$  is the transition guard function.
- $A:(Q\times Q)\to (X\to X)$  is the transition action.

Our systems are autonomous, transition time points are measured relative to the start of the system, all differential equations are autonomous. In the following we use  $q_{i_1}, q_{i_2}, \cdots$  to denote a sequence of modes with indices  $i_1, i_2, \cdots$ . An *execution*  $\epsilon$  of the hybrid system HS starting at some time point  $t_0^1$ , mode  $q_0$  and state  $x_0$  is a sequence

$$\epsilon = (q_0, t_0, \Phi_0), (q_{i_1}, t_1, \Phi_1), \cdots$$

such that  $(q_{i_j},q_{i_{j+1}})\in \delta$ ,  $\Phi_j(t)$  is the solution of the equations  $F(q_{i_j})$  for  $t_j\leq t\leq t_{j+1}$ , with initial conditions  $\Phi_j(t_j)=\Phi_{j-1}(t_j), j>0$ , and  $\Phi_0(t_0)=x_0$ .

<sup>&</sup>lt;sup>1</sup>This time is relative to the start of the system, the system itself is autonomous

Further it must hold that  $I(q_i)(\Phi_{t_j}(t)) = \text{true}$  for  $t_j \leq t < t_{i+j}, \ G(q_{i_j}, q_{i_{j+1}})(\Phi_j(t_j)) = \text{true}$ , and  $A(q_{i_j}, q_{i_{j+1}})(\Phi_j(t_j)) = \Phi_{j+1}(t_j)$ .

A trace  $\tau$  starting at time  $t_0$  in mode  $q_0$  and state  $x_0$  is a (finite or infinite) sequence of modes together with the time points at which the system switched into that mode:

$$\tau = (q_0, t_0), (q_{i_1}, t_1), (q_{i_2}, t_2), \cdots$$

such that there is a unique execution  $\epsilon$  with the same modes and times.

As noted, the hybrid systems 2.1 are non-deterministic. This means that the set of all traces starting in some mode  $q_0$ , state  $x_0$ , and time  $t_0$ , denoted with  $T_{q_0,x_0,t_0}$ , can contain more then one element, it is in general infinite.

There could be traces that end in a blocking state, meaning a mode q and state x where I(q)(x)= false and  $\forall \tilde{q}.\ (q,\tilde{q})\in \delta \to G(q,\tilde{q})=$  false. Also, there could be so called zeno traces, traces where  $\lim_{i\to\infty} t_{i+1}-t_i\to 0$ .

A transition is called a *forced transitions* if it happens at a time point t in some mode q and state x where I(q)(x(t)) is false and G(q,q')(x(t))= true.

We call a hybrid system *proper* if all forced transitions are unique: if the system is in the mode q, state x at time point t and I(q)(x(t)) = false then there exists at most one q' such that G(q,q')(x(t)) = true.

**Definition 2.2 (Event)** Given a hybrid system HS and a set of traces  $T_{q_0,x_0,t_0}$ . An event  $e_{q_i,q_j,t}$  is a mapping from the set of traces to the set of traces that selects all the traces  $\tau_i \in T_{q_0,x_0,t_0}$  containing the transition from  $q_i$  to  $q_j$  at time point t. the set of all traces that change mode from  $q_i$  to  $q_j$  at time point t

Note,  $e_{q_i,q_j,t}(T_{q_0,x_0,t_0})$  could be the empty set.

Finally, we can define the notion of a controlled trace, being the set of traces in which all transitions are either forced of are caused by an event.

**Definition 2.3 (Controlled Traces)** To a sequence of events  $< e_{q_{i_j},q_{i_{j+1}},t_j} >$ , the set of traces  $\tau$  with  $\tau \in T_{q_0,x_0,t_0}$ , if they exists, that contain only forced transitions and transitions caused by the events is called a controlled traces.

**Lemma 2.4** For a proper hybrid system HS, a sequence of events and a set of traces the controlled trace, if exists, is unique.

The proof consists of applying the definitions. All forced transitions are deterministic by the properness assumption, all other transitions in a controlled trace correspond to events.

### 3 THE CONTROL PROBLEM

Equipped with hybrid systems we now model the control system by combining the model of the plant, constraints of physical and logical nature, and the possible control inputs, given in form of control functions, into one hybrid system.

The system to be controlled is represented by the system of differential equations

$$\dot{x} = f(x, u) \tag{1}$$

where  $x \in \mathbb{R}^n$ ,  $u \in \mathbb{R}^m$ ,  $f : \mathbb{R}^n \times \mathbb{R}^m \to \mathbb{R}^n$  models the dynamics of the plant. We assume that the plant is totally observable and hence, for simplicity, do not have an explicit output function. Further the system is assumed to be controllable and  $f(\mathbf{0},\mathbf{0}) = \mathbf{0}$ , and f is Lipschitz continuous with constant L

$$||f(x,u) - f(\tilde{x},u)|| \le L||x - \tilde{x}||.$$

The control input u can not be chosen freely, but the value of u is computed by a member of a set of control functions  $\operatorname{Cf}_i, i=1\cdots k$ ,  $\operatorname{Cf}:\mathbb{R}^n\to\mathbb{R}^m$ , (Morse, 1993), the control algorithm plays the role of a supervisor and determines the time points when to use which controller. For one of the  $\operatorname{Cf}_i$  it holds that  $\operatorname{Cf}_i(0)=0$ . This is motivated by systems in which digitally controlled actuators interact with the plant.

The state and the control input are constrained. The constraints are given in form of a system of equalities and inequalities. We do not restrict to linear constraints, since we intend to solve a non-convex problem. To model the movements of a car in a maze we need to include the dimension of the car, collision with walls is determined by deciding if a given set of equalities and inequalities does not have a solution. There could also exist logical constraints, constraints that dictate sequencing requirements such as  $\dot{u}$  should be zero before u can change from positive to negate, e.g. the car can not reverse before it came to a complete stop.

The goal of the control algorithm is to minimize the cost function

$$J_t(x,u) = \int_0^t g(x,u)dt$$

for a positive definite function g. It holds that  $g(\mathbf{0},\mathbf{0})=0$ . The control algorithm is restricted in that it does not have a global view of all constraints, but can only see constraints that are "visible" within a given *control horizon H*. This horizon can not be chosen freely, it must by sufficiently large such that the resulting recent horizon controller is stabilizing (Almir and Bornard, 1995). We examine this further in 4.1.1.

The control problem is modeled using a hybrid system with k modes. (k is the number of control functions.) In each mode the equations  $F(q_i)$  are given

by the model equation  $\dot{x} = f(x, u)$  and the equation  $u = cf_i(x)$ . The mode invariants model the constraints. The transitions encode the sequencing requirements and the transitions are guarded by the mode invariant they lead to. This is to ensure that the system does not switch into a mode in which the invariant is false. Note, for technical reasons, each mode  $q_i$  contains a self transition  $\delta(q_i, q_i)$  guarded by the invariant of the mode. To ensure that the control system remains at least a minimal amount of time  $t_{\min}$ , minimal holding time, in each mode we add the state variable, clock, with the differential equation clock = 1 to each mode, each transition guard is augmented with  $clock>t_{\min}$  and  $t_{\min}$  is set to 0 during each transition by A. The developer has to ensure that the constructed hybrid system is proper by determining which transition should be taken in case a mode invariant becomes false and modify the transition guards accordingly.

The resulting hybrid system is called the *controlled hybrid system* CHS. The control problem hence is reduced to "switching the mode of the CHS system at runtime".

For convenience, for any trace  $\tau \in T_{q,x,t}$  we write  $J(\tau)$  to denote the cost J computed over the trace  $\tau$ . The value of  $J(\tau)$  can be computed over each trace since x and u are state variables.

The control problem can not be solved in general. We need assumptions under which an algorithm will succeed.

**Assumption 1** For  $x_0$  and  $q_0$  the CHS is non-blocking, there exists an trace  $\tau$  of infinite length.

For a system CHS, the set of the non blocking traces is denoted with NBT $^{\infty}$ . The infinity stands for the fact that the set NBT $^{\infty}$  could be constructed by an infinite construction. The optimal solution of the control problem is the trace  $\tau \in \text{NBT}^{\infty}$  with minimal  $J(\tau)$  denoted as  $\tau^*$ .

**Definition 3.1 (CHS Control Problem)** Given a hybrid system CHS, an initial state  $x_0$ , initial mode  $q_0$  and the cost function J, if  $\mathsf{NBT}^\infty$  is not empty then the sequence of events  $< e_{q_{i_j},q_{i_{j+1}},t_j} >$  that corresponds to  $\tau^*$  is the optimal solution of the CHS control problem

By construction the control hybrid system CHS remains for at least  $t_{\min}$  time units in each mode.

## 4 THE CONTROL ALGORITHM

For discrete systems, recent horizon control algorithms approach the problem by repeatedly solving a finite horizon control problem, realizing the first control move and then repeating the procedure. We

construct an equivalent process for our mixed discrete continuous system.

As noted, the control algorithm can not explore the entire future behavior of the system but is restricted to the prediction horizon  $\cal H$ .

**Definition 4.1** (*H*-**Predictive Trace**) Let CHS be a hybrid system. A *H*-predictive trace starting in mode q, and state x and time point t is a trace  $\tau = (q,t), (q_{i_1},t_1), \cdots (q_{i_n},t_n)$  of the CHS hybrid system starting in mode q, state x and has duration H,  $t_n = t + H$ . The set of all these traces is denoted with  $T_{q,x,t}^H$ .

In general the set  $T_{q,x,t}^H$  is infinite. In practice we have to impose restrictions such that we can explore this set. We use two constants: (1) The previously defined  $t_{\min}$ , which is the minimal amount of time the system has to remain in each of the modes in any trace. The reason to restrict the minimal amount of time the system has to remain in a mode to  $t_{\min}$  is to gives us the time needed to run the control algorithm. (2)  $t_{\max}$ , which is the maximal amount of time the system can remain in any mode without considering to switch mode.  $t_{\min},t_{\max}$  gives us the time frame in which control moves are scheduled to happen. These restrictions now allow us to reduce the size of  $T_{q,x,t}^H$ :

**Definition 4.2 (Restricted** *H***-Predictive)** *A* restricted *H*-predictive trace starting in mode q and state x is a predictive trace

$$\tau = (q, t), (q_{i_1}, t_1), (q_{i_2}, t_2), \cdots (q_{i_n}, t_n)$$

where  $t_{k+1} - t_k$ , the time spend in mode  $q_k$ , is uniquely determined as the maximal amount of time the system can remain in mode  $q_{i_k}$  while the invariant  $I_{q_{i_k}}$  holds and  $t_{min} \leq t_{k+1} - t_k \leq t_{max}$ .

By construction we have  $t_{k+1} - t_k \ge t_{\min}$ .

Note, the CHS hybrid system contains self transitions for each mode, this means in a predictive trace the system could remain in any mode as long as the invariant of this mode is not violated. We use  $\operatorname{RT}_{q,x}^H$  to denote the set of all restricted H-predictive traces that start in mode q and state x By construction and lemma 2.4,  $\operatorname{RT}_{q,x}^H$  is a finite set. Transitions either correspond to events or are forced. The length n of the H predictive traces is at most  $\lfloor H/t_{\min} \rfloor$  and the at least  $\lfloor H/t_{\max} \rfloor$ .

Now we are ready to give the control algorithm:

**Definition 4.3 (Supervisor)** Let  $t_i$  be the next time point where a control move is required, and  $t_{run}$  be a constant given at the end of this definition.

1. At time point  $t_i$ — $t_{run}$  the supervisor reads the mode q and state x of the system and determines all restricted H-predictive traces  $RT_{x,a}^H$ :

$$\tau = (q, t_i - t_{run}), (q_{i_1}, t_i), (q_{i_2}, t_2), \cdots (q_{i_n}, t_n)$$

which are the traces that start in mode q remain for  $t_{run}$  in mode q and then switch to some new mode  $q_{i_1}$ , which could be q again.

- 2. If  $RT_{x,q}^H$  is the empty set, meaning there is no control move that prevents the violation of the constraints the system has to be stopped or invoke a user defined ad hoc procedure to relax constraints so that the system becomes feasible again. Let  $\tau$  be the trace with minimal  $J(\tau)$  of the traces  $\tau \in RT_{x,q}^H$ .
- 3. At time point  $t_i$  the supervisor causes the hybrid system to switch to the mode  $q_{i_1}$ .
- 4. The supervisor waits until the time point  $t_2 t_{run}$  and starts again with step 1.

 $t_{run}$  is the amount of time the supervisor needs to perform its computation.

The algorithm should also keep track of the value of  $J(\tau)$  computed in the previous step and check that it is larger then  $J(\tau)$  computed in the current step which allows to verify that the plant progresses towards the minimum of J as detailed in the next section.

### 4.1 Stability and Complexity

There is a documented gap between the theoretical stability results for recent horizon controllers and there practical use in industrial applications (Allgöwer and Findeisen, 2002; Morari and Lee, 1999; Davrazos and Koussoulas, 2001). Stability is assured by either introducing carefully chosen terminal weights, defining a set that has to be reached within each of the finite predictions, or by determining the prediction horizon such that repeatedly solving the finite horizon problem results in a global optimal solution (Mayne and Michalska, 1990; Primbs and Nevistic, 2000; Almir and Bornard, 1995). We relay on the latter approach by determining assumptions under which the proposed algorithm is stabilizing.

#### 4.1.1 Stability

To show stability of the proposed control algorithm we relay on the classic Liapunov argument using the costs  $J_H(x,q)$  computed over the finite horizon H, which are the predictive traces, as Liapunov function (Mayne and Michalska, 1990; Primbs and Nevistic, 2000; Almir and Bornard, 1995). We have to determine under which conditions the sequence of the  $J_H(x,q)$  is decreasing.

To ensure convergence we assume that the prediction horizon is sufficiently large such that the sequence of the costs  $J(\tau_i)$  where  $\tau_i$  is the predictive trace computed in the *i*-th step of the control algorithm is decreasing.

**Assumption 2** For all traces

$$\tau = (q_0, t_0), (q_{i_1}, t_1), (q_{i_2}, t_2), \cdots$$

starting in state  $x_0$  and mode  $q_0$ , with  $\tau \in \mathsf{NBT}^\infty$ , there is an H such that for all the finite traces  $\tilde{\tau} \in \mathsf{RT}^H_{q_{j_i},x_i}$  and  $\overline{\tau} \in \mathsf{RT}^H_{q_{j_{i+1}},x_{i+1}}$ ,  $\tilde{\tau}$  and  $\overline{\tau}$  being the optimal traces, where  $x_i$  is the i-the state and  $q_{j_i}$  the i-th mode of  $\tau$ , there is an  $\alpha$ ,  $0 < \alpha < 1$  such that

$$\alpha J(\tilde{\tau}) \leq J(\overline{\tau}).$$

Assumption 2 ensures that the finite horizon strategy can succeed in solving the control problem by requiring that H was choose such that for all predictive traces  $\tau_i$ 

$$\sup \frac{J(\tau_i(i+1))}{J(\tau_i)} < 1.$$

This represents the conditions of the stability criteria as found in the above mentioned literature.

Let  $\tau$  be the trace and  $\epsilon$  be the corresponding execution. To apply the stability criteria we translate our continuous system into a discrete sequence of states by using the solutions  $\Phi_i$  of the execution, hence

$$x(n+1) = \Phi_i(x(n), t_{n+1})$$

This sequence is well defined since the solution to the initial problems as we enter mode  $q_{i_n}$  is unique.

**Theorem 4.4 (Stability)** *Under the assumptions 1,2 the control algorithm is stable.* 

We have to show that repeatedly computing the predictive traces, determining a sequence of events we obtain a trace that approximates the optimal solution  $\tau^*$  of the control problem. Let  $\tau_i$  be the sequence of the predictive traces choose by the control algorithm.

Let x(j) and  $q_{i_j}$  be the state and mode we are in at time point  $t_j$ , j=0 is the initial state.  $q_{i_{j+1}}$  is determined from the trace  $\tilde{\tau} \in \mathrm{RT}^H_{q_{j_i},x(i)}$  with minimal cost. The system moves to mode  $q_{i_{j+1}}$  at time point  $t_{k+1}$  where we are in state x(j+1). Let  $\overline{\tau}$  be the next traces determined by the algorithm 4.3. Be assumption 2 we have that either that  $J(\overline{\tau}) < J(\tilde{\tau})$  since  $\alpha < 1$ , or they are equal in the case that we already reached the stable point x=0. By this the cost of the local finite horizon predictions as computed by the algorithm  $J(\tau_i)$  is decreasing and with this the system is stable.

#### 4.1.2 Complexity

The complexity of our approach is determined by the size of the set  $\mathrm{RT}_{q,x,t}^H$  which grows exponentially in the number of mode switches of the predictive traces. However, in practical applications the situation is not as bad as it first seems. If there is no "constraint visible" within the prediction horizon the predictor uses

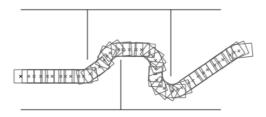


Figure 2: Track vehicle

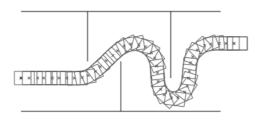


Figure 3: Wheeled vehicle

large steps, defined by  $t_{max}$ , to reach the prediction horizon. If there are constraints then traces resulting in a blocked state are cut reducing the search space. The observation of the reduction in search space caused by cutting branches of the search tree is consistent with observations made by other approaches to model predictive control (Abdelwahed et al., 2002), (Hadj-Alouane et al., 1994). Also, our implementation uses symbolic techniques based on (v. Mohrenschildt, 2001), which we do not detail here, that reduce the amount of computation needed at run-time by symbolically pre-computing the solution of the differential equations whenever possible. The symbolic solutions allow us to reduce the real- time computation to evaluation of formulas. Obviously, the number of formulas is still exponential with is consistent with the explicit MPC algorithm algorithm presented in (Bemporad et al., 2002).

# 5 NONHOLONOMIC CAR-LIKE ROBOT

We used our approach to control a nonholonomic robot moving in a maze. We present both, a car with tracks and a car with wheels. For details about the modeling and control of nonholonomic cars the reader is referred to the book (Laumond, 1998). The example presented here was inspired by (Almeida et al., 1997).

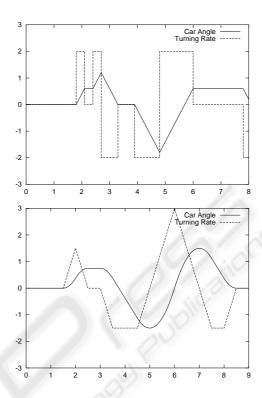


Figure 4: Control input and angle for tracked and wheeled vehicle

First we like to point out that in contrast to many approaches to car control there is no path planing, meaning that the car is not tracking a predefined path. The controller can only explore the possible paths within the prediction horizon and "sees" the walls only as they enter into the prediction horizon. Also, in contrast to (Almeida et al., 1997) we did not develop a special controller for this problem but run the controller as defined in this paper in its generic form.

The example demonstrates that our control algorithm is able to perform point-to-point motion control without path planing. The prediction horizon H was chosen such that the controller was able to "look around the corner", which means, as we move with 1 m/sec and the longest wall is 1m long we use a prediction horizon of H=1sec.

The car is modeled as follows:

$$\dot{x} = v \cos(\theta), \quad \dot{y} = v \sin(\theta), \quad \dot{\theta} = vu \quad (2)$$

with the control variable u and  $v \in \{-1,0,1\}$ . x,y is the position of the center of the rear steering axle,  $\theta$  is the rotation of the car in relation to the coordinate system. The four sides of the car are computed at runtime using  $x,y,\theta$ , using the the car length .7, width .4, and offset .2 of the rear axle.

We performed two types of simulations, u being discrete valued,  $u \in \{-2, 0, 2\}$ , which models the

movement of a track vehicle, and u being change rate constrained,  $u \in [-3,3]$ ,  $du \in \{-5,0,5\}$ , which models a car more realistically. We extended the constraints for u for the wheeled car from 2 to 3 since else the car was not able to turn sufficiently fast.

For both simulations we used a mixed numeric, symbolic approach. To detect if the car is hitting a wall, which means solving 4 linear systems, we generated symbolic code using the computer algebra system Maple. We also solved the model equations symbolically, but to predict the time which we can remain in a mode, we use numeric methods. We control the car to move to the point (7, 2), which is at the right side middle of the maze in the figures 2.

The Figures in 2 show the movement of the car. The first figure plots the movement of a track vehicle, second of a car like vehicle. Figure 4 plots the control input being the change of the turning rate of the car versus the angle of the car again for the discrete control input and for the change rate limited control input.

#### 6 CONCLUSIONS

We presented an approach to model the influence of discrete control moves using hybrid systems and presented a control algorithm. In contrast to other approaches we use a continuous model of the plant, the continuous control input is computed my a member of the control functions. The presented algorithm is adaptive in the sense that the control frequency, the frequency at which control actions occur is not fixed but changes with the presence of constraints.

The controller was implemented for several case studies, not limited to the here presented maze tracking problem, and we found that we could solve complex control tasks using our generic controller.

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